

# Information efficiency in investment

Thomas M. Cover and Elza Erkip<sup>1</sup>

Stanford University, Information Systems Lab, Stanford, CA 94305-4055  
 cover@isl.stanford.edu, elza@isl.stanford.edu

*Abstract* — We answer the question, what should we say about  $V$  when we want to gamble on  $X$ , and what is it worth? If  $V = X$ , we show that every bit of description at rate  $R$  is worth a bit of increase  $\Delta(R)$  in the doubling rate. Thus the efficiency  $\Delta(R)/R$  is equal to 1. For general  $V$ , we provide a single letter characterization for  $\Delta(R)$ . When applied specifically to jointly normal  $(V, X)$  with correlation  $\rho$ , we find the initial efficiency  $\Delta'(0)$  is  $\rho^2$ . If  $V$  and  $X$  are Bernoulli random variables connected by a binary symmetric channel with parameter  $p$ , the initial efficiency is  $(1 - 2p)^2$ .

We finally show how much increase in doubling rate is possible when the sender can provide  $R$  bits of information about  $V$  and side information  $S$  is available only to the investor.

## SUMMARY

Suppose we are interested in gambling on the outcome of a random variable  $X$ . The gambling consists of betting a proportion of wealth  $b(x)$  on the outcome  $x$ . We would like to maximize the doubling rate, which is the growth rate of wealth when the gambler uses a fixed betting strategy on independent realizations of  $X$ . It is well known that Kelly gambling, which is to bet in proportion to the probability mass function of  $X$ , is optimal.

Now suppose we allow a description of  $X$  at rate  $R$  bits per symbol. Let  $\Delta(R)$  be the maximum increase in the doubling rate of wealth for transmission rate of  $R$ . We prove that  $\Delta(R) = R$ . Any bit of information one sends about  $X$  is worth a bit of increase in the doubling rate.

We next consider the effectiveness of sending information when side information  $S$  is available to the investor but not to the encoder. The gambler combines this side information with the partial description of  $X$  to form the bet.

**Theorem 1** *If  $X$  is described at rate  $R$ , and side information  $S$  is available to the gambler, then,*

$$\Delta(R) = R.$$

We ask what information should be given about a correlated random variable  $V$  if we want to help the investor gamble on  $X$ . This problem shows some similarities to source coding with side information [4, 1]. The encoder sends  $R$  bits about  $V$  and the investor uses this information to gamble on  $X$ . Here maximal efficiency is not generally possible.

**Theorem 2** *When the encoder observes  $V$  correlated with  $X$ ,*

$$\Delta(R) = \max_{p(\tilde{v}|v,x): I(\tilde{V};V) \leq R, \tilde{V} \rightarrow V \rightarrow X} I(\tilde{V}; X).$$

We establish certain properties of  $\Delta(R)$  using entropy maximization results from Witsenhausen and Wyner [3].

Next, we find the increase in the doubling rate when the encoder sends information at rate  $R$  about a correlated random variable  $V$  with side information  $S$  present only at the investor. The investor uses these  $R$  bits together with the side information  $S$  to invest in the outcome of  $X$ .

**Theorem 3** *When the encoder observes  $V$ , and side information  $S$  is available at the investor,*

$$\Delta(R) = \max_{p(\tilde{v}|v,x,s): I(V;\tilde{V}|S) \leq R, \tilde{V} \rightarrow V \rightarrow (X,S)} I(\tilde{V}; X|S)$$

Finally, we investigate the efficiency of descriptions based on correlated variables. If  $X$  and  $V$  are both Bernoulli( $\frac{1}{2}$ ) and are associated by a binary symmetric channel with crossover probability  $p$ , it can be shown that  $\Delta(R)$  has a derivative of  $(1 - 2p)^2$  at  $R = 0$ . Thus, even the most effective description of  $V$  relative to the investment in  $X$  pays off at the rate of only  $(1 - 2p)^2$  bits of doubling per bit of description.

Now suppose that  $V$  and  $X$  are jointly Gaussian with correlation  $\rho$ . In this case the initial efficiency,  $\Delta'(0)$ , is equal to  $\rho^2$ .

The functional form of  $\Delta(R)$  for binary and Gaussian random variables will be developed in [2]. Also, the relationship between the derivative of  $\Delta(R)$  at  $R = 0$  and the Renyi maximal correlation of  $V$  and  $X$  will be investigated.

## REFERENCES

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